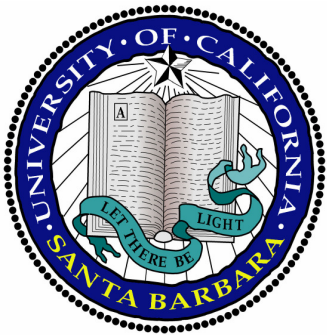




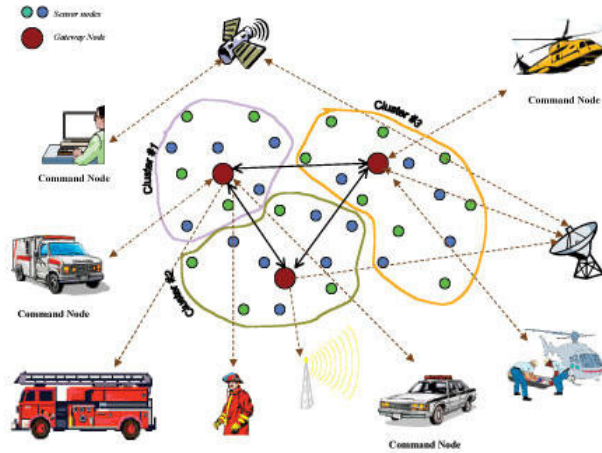
# Code Design for Fast Selective Retrieval of Fusion Stored Sensor Network/Time- Series Data

Sharadh Ramaswamy, Jayanth Nayak and Kenneth Rose

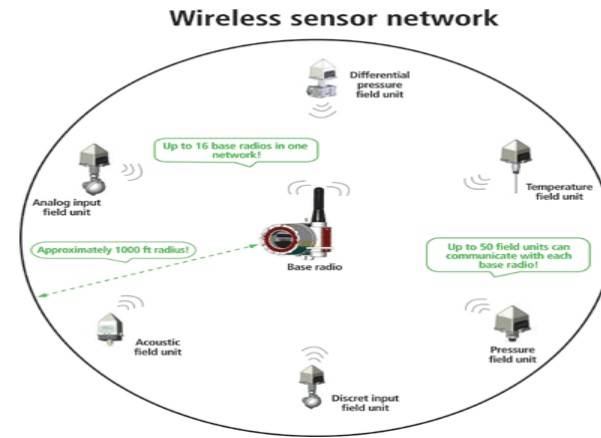


# Monitoring with Sensor N/Ws

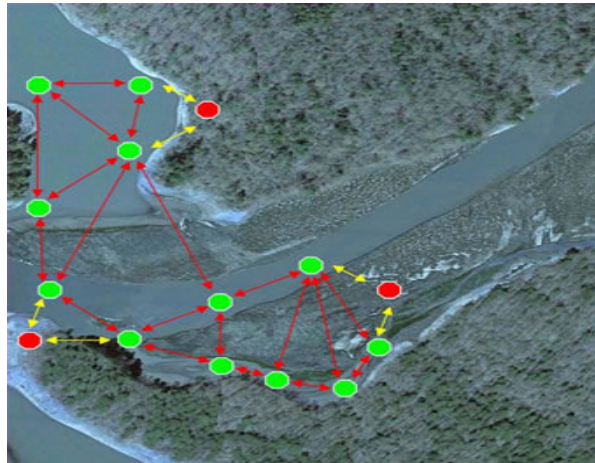
<http://faculty.cua.edu/elsharkawy/images/cluster.jpg>



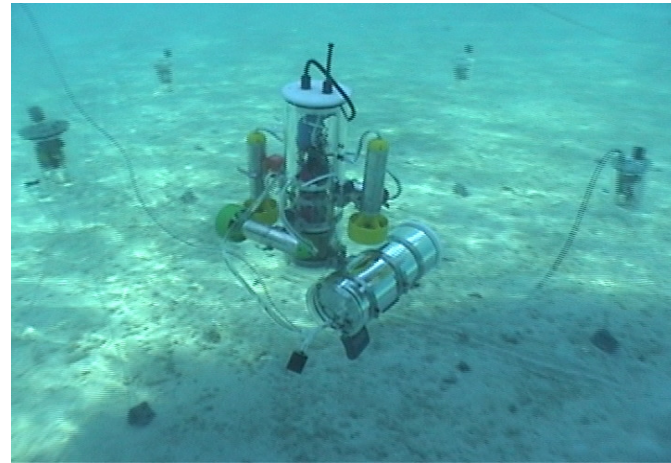
<http://www.isa.org/Images/IntTech/2005/July/2005070553-3.gif>



[http://gis.clemson.edu/cpost/images/sensor\\_network\\_small.jpg](http://gis.clemson.edu/cpost/images/sensor_network_small.jpg)



[http://groups.csail.mit.edu/dri/wiki/images/d/db/amo\\_ur\\_with\\_sensors.jpg](http://groups.csail.mit.edu/dri/wiki/images/d/db/amo_ur_with_sensors.jpg)

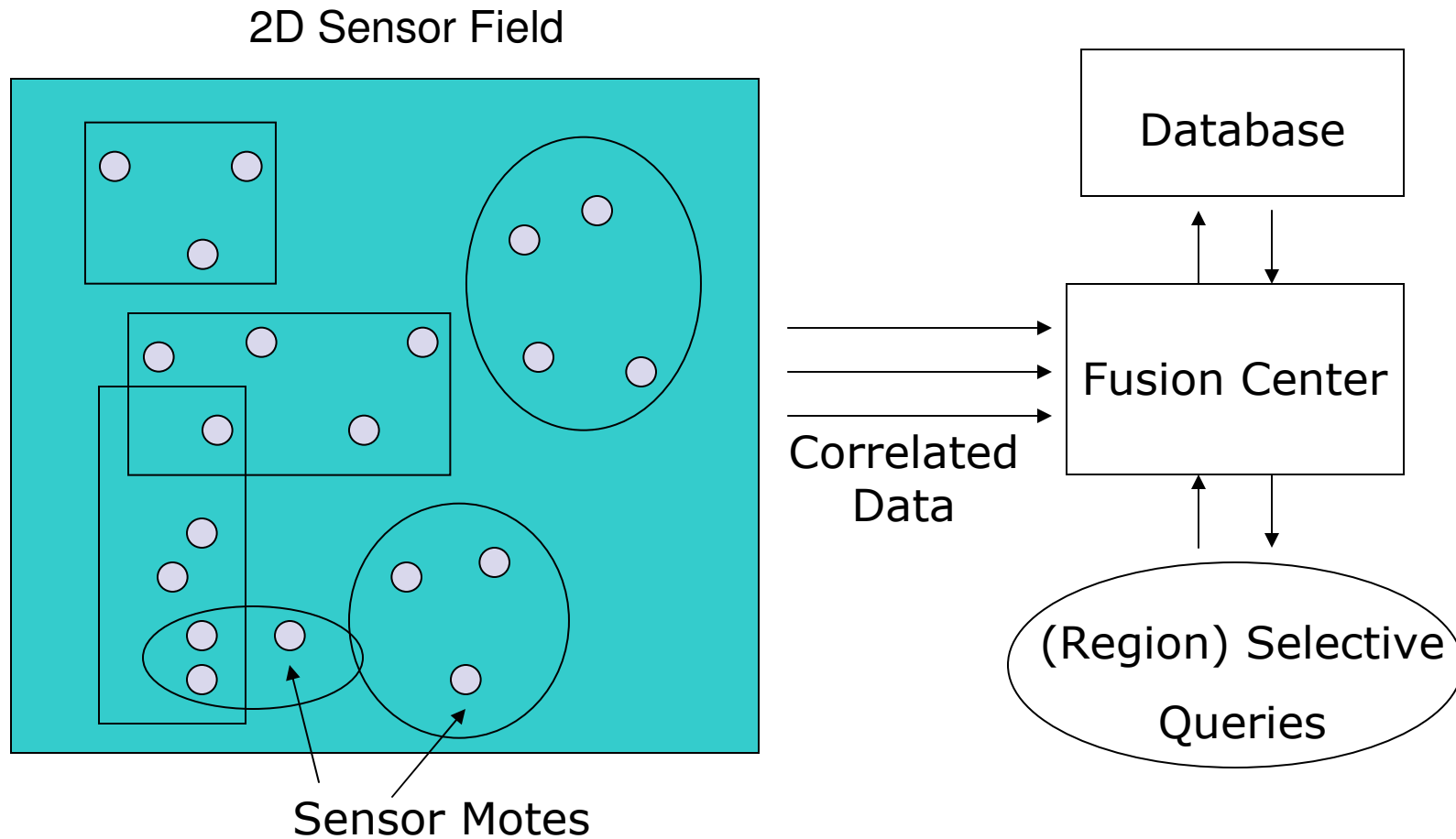


# Motivation

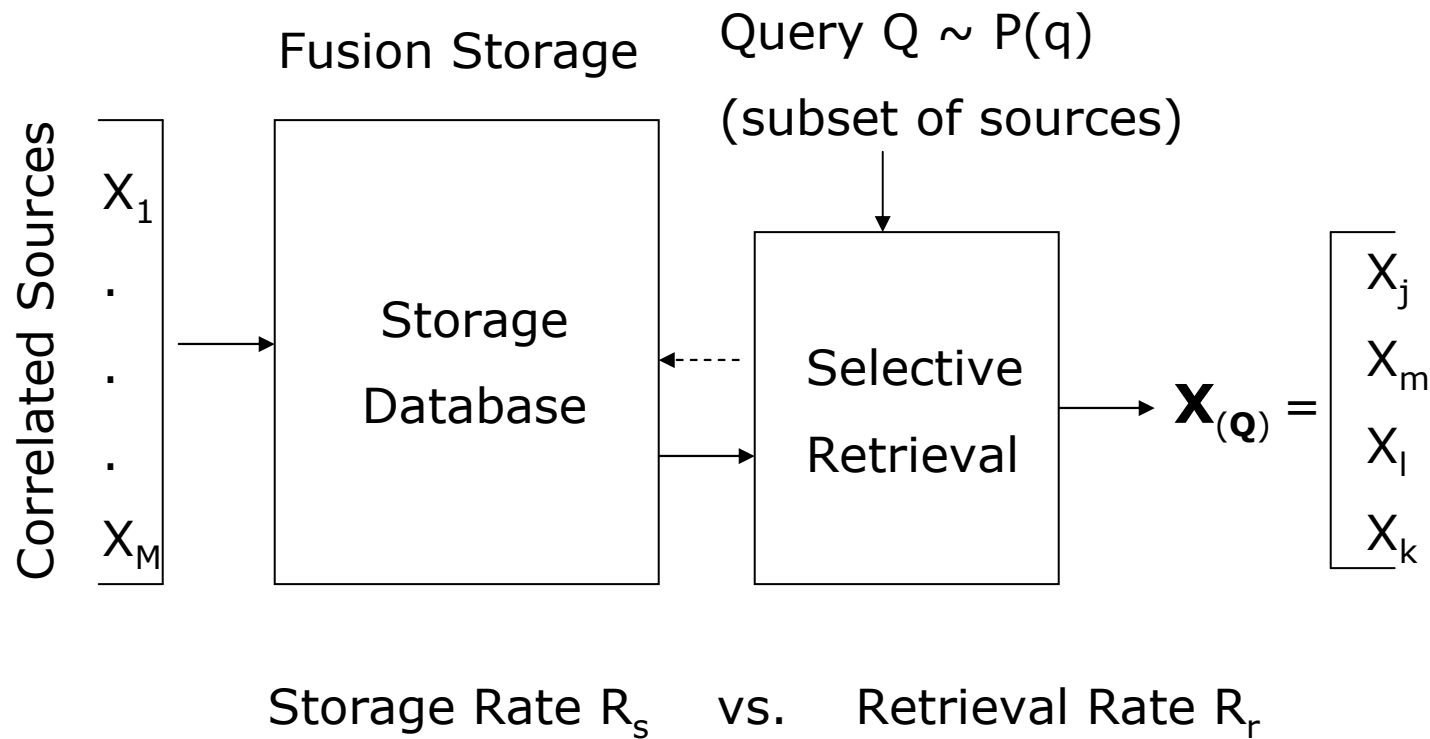


- Dense sensor N/Ws in many apps. (monitoring/tracking/surveillance)
- Correlated measurements are collected and stored
- Distributed compression for efficient data collection
- **Optimal storage strategy?**

# Typical User/Query Model



# Fusion Storage vs. Selective Retrieval



# Fusion Storage vs. Selective Retrieval

- Fundamentally a *storage* problem
- Challenge: exploit correlations to
  - min. storage rate
  - min. retrieval rate
- Conflicting objectives ?!
- Trade-off characterized in prior work (Nayak et. al. ISIT'05)

# Naive Storage Strategy I

- **Jointly compress all sources**

- optimal in exploiting inter-source corr.

- **minimizes storage rate**  $R_s$

Lossless Coding  $\Rightarrow R_{s,\min} = H(X_1, \dots, X_M)$

- retrieves all stored data even for a small subset

- **high retrieval rate/time!!**  $R_r = R_s$

# Naive Storage Strategy II

- **Compress and store every subset of sources separately**

- retrieves only minimum info. reqd.
- **optimal in retrieval rate/time**

$$\text{Lossless} \Rightarrow R_{r,\min} = \sum_{\mathbf{q}} P(\mathbf{q}) H(\mathbf{X}_{(\mathbf{q})})$$

- reqd. storage grows with size of query set
- **(combinatorially) high storage rate!!**

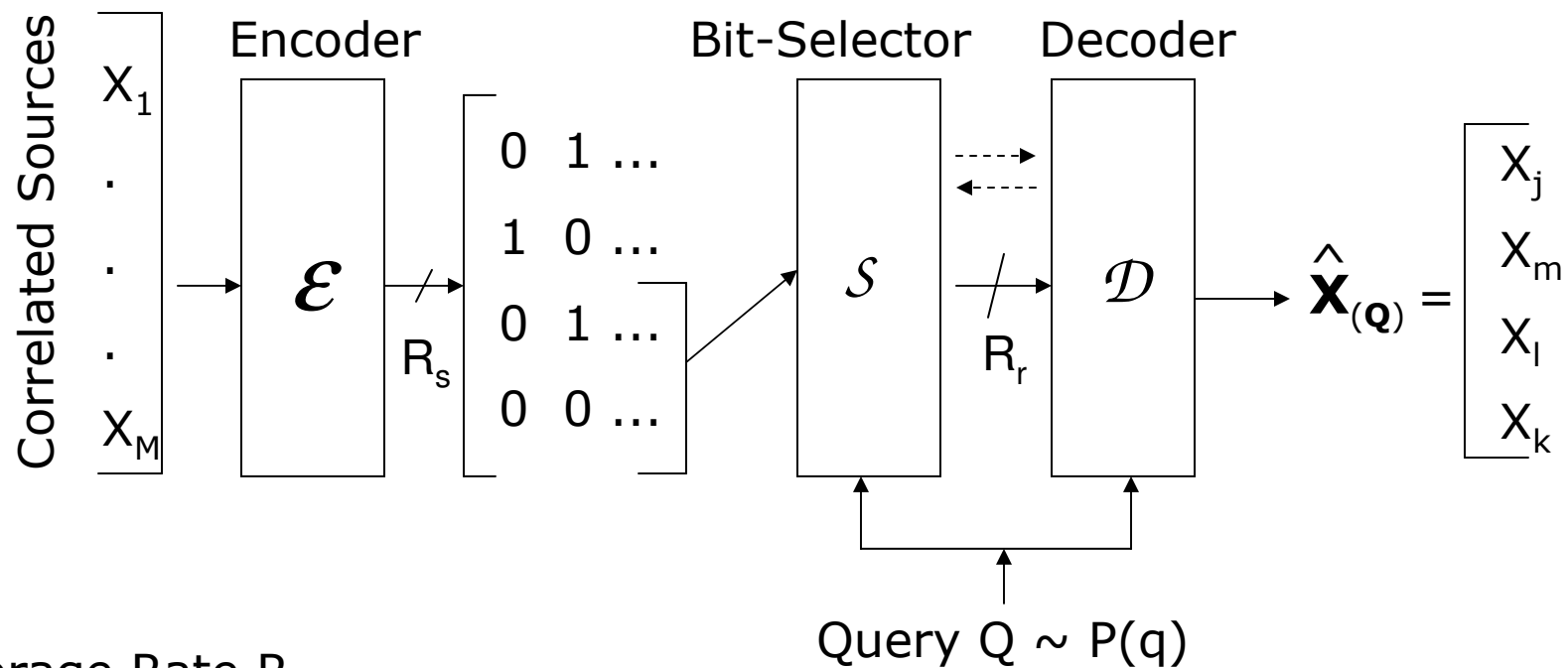
$$\text{Lossless} \Rightarrow R_s = \sum_{\mathbf{q}} H(\mathbf{X}_{(\mathbf{q})}) \gg H(X_1, \dots, X_M)$$



# Storage with Distortion

- Perfect reconstruction not always necessary
- Trade-off -
  - Storage Rate vs. Retrieval Rate vs. Distortion
- In practice, storage devices are of fixed capacity
- Practical design by fixing maximum allowable  $R_s$
- Query-dependent “bit-selection” (and relevant codebooks) for selective retrieval...

# Proposed Fusion Coding Approach



Storage Rate  $R_s$   
 Distortion  $D$   
 Retrieval Rate  $R_r$

# Mathematically...

$$\text{Encoder: } \mathcal{E} : \mathcal{R}^M \rightarrow \mathcal{I} = \{0, 1\}^{R_s}$$

$$\text{Decoder: } \mathcal{D} : \mathcal{I} \times \mathcal{B} \rightarrow \hat{\mathcal{X}} \quad \text{Bit-Selector: } \mathcal{S} : \mathcal{Q} \rightarrow \mathcal{B} = 2^{\{1, \dots, R_s\}}$$

$$D = \sum_{\mathbf{q} \in \mathcal{Q}} P(\mathbf{q}) \frac{1}{|\mathcal{X}|} \sum_{\mathbf{x} \in \mathcal{X}} d_{\mathbf{q}}(\mathbf{x}, \hat{\mathbf{x}}) \quad R_r = \sum_{\mathbf{q}} P(\mathbf{q}) R_{\mathbf{q}} = \sum_{\mathbf{q}} P(\mathbf{q}) |\mathcal{S}(\mathbf{q})|$$

$$\text{Objective: } \min_{\mathcal{E}, \mathcal{S}, \mathcal{D}} J = D(R_s) + \lambda R_r(R_s), \lambda \geq 0$$

# Optimality and Design

Optimal Encoder:  $\mathcal{E}(\mathbf{x}) = \arg \min_{\mathbf{i} \in \mathcal{I}} \sum_{\mathbf{q}} P(\mathbf{q}) d_{\mathbf{q}}(\mathbf{x}, \mathcal{D}(\mathbf{i}, \mathcal{S}(\mathbf{q}))), \forall \mathbf{x}$

Optimal Bit-Selector:  $\mathcal{S}(\mathbf{q}) = \arg \min_{\mathbf{e} \in \mathcal{B}} \left\{ \frac{1}{|\mathcal{X}|} \sum_{\mathbf{x}} d_{\mathbf{q}}(\mathbf{x}, \mathcal{D}(\mathcal{E}(\mathbf{x}), \mathbf{e})) + \lambda |\mathbf{e}| \right\}, \forall \mathbf{q}$

Optimal Decoder:  $\mathcal{D}(\mathbf{i}, \mathbf{e}) = \frac{1}{|F|} \sum_{\mathbf{x} \in F} \mathbf{x}, \forall \mathbf{e}, \mathbf{i}$

Design by  
Gradient  
Descent

where  $F = \{\mathbf{x} : (\mathcal{E}(\mathbf{x}))_{\mathbf{e}} = (\mathbf{i})_{\mathbf{e}}\}$ .

# Experimental Set-up

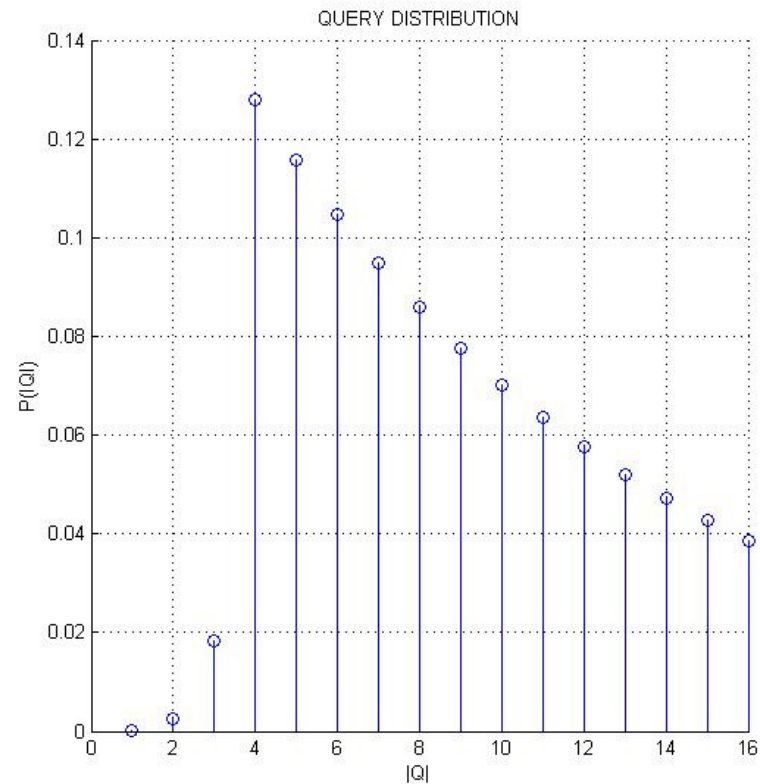
- Sensor data modeled as corr. Gaussian

$$C = \sigma^2 \begin{pmatrix} 1 & \rho & \dots & \rho^{M-2} & \rho^{M-1} \\ \rho & 1 & \rho & \dots & \rho^{M-2} \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ \rho^{M-2} & \dots & \rho & 1 & \rho \\ \rho^{M-1} & \rho^{M-2} & \dots & \rho & 1 \end{pmatrix}$$

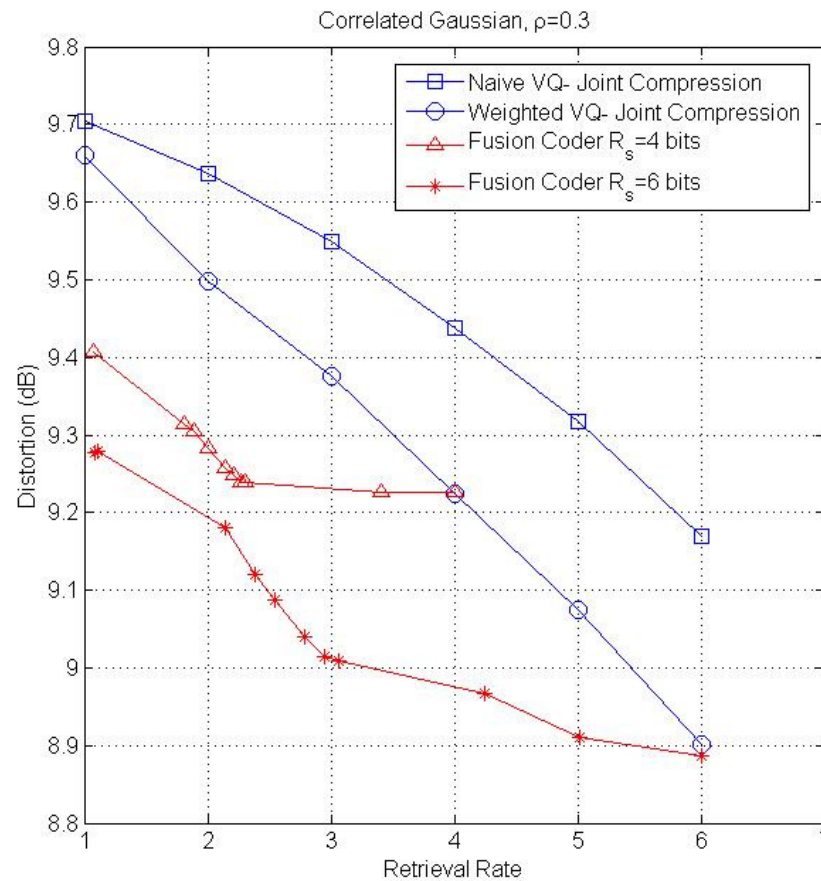
- Stock market data (UCR Data-mining archive) – 93 stocks

# Exponential Query Distribution

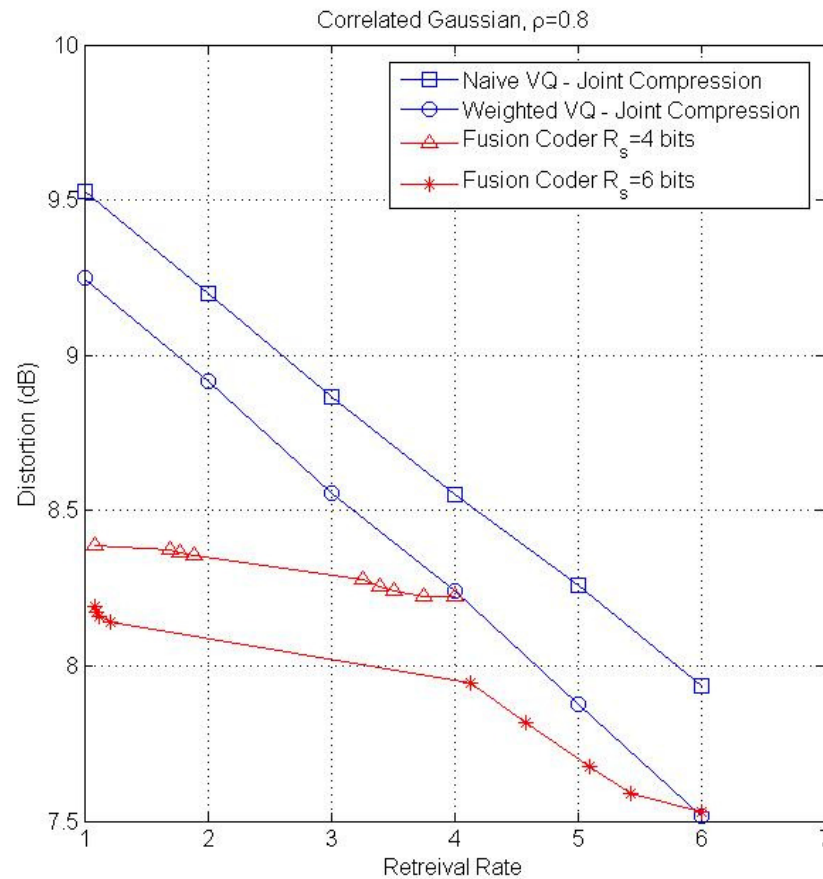
- Query distribution modeled as exponential in query size
- Queries of same size equally likely
- Distribution approximated by a training set of 335 queries



# Correlated Gaussian $\rho=0.3$

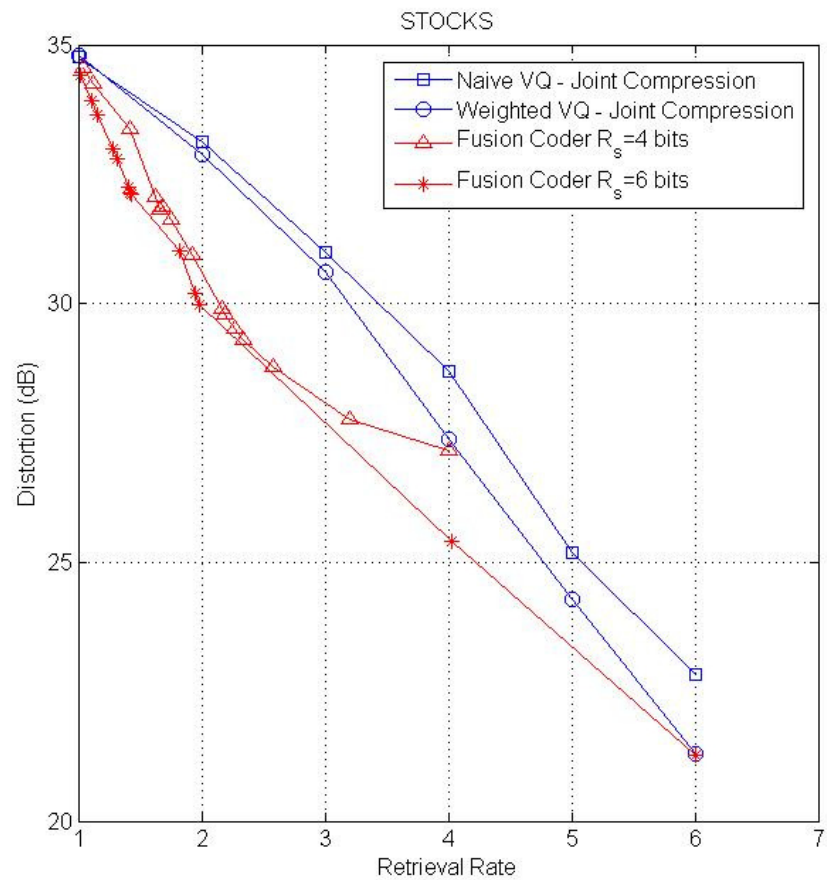


# Correlated Gaussian $\rho=0.8$





# Stock Market Data





# Conclusions

- Fusion storage vs. selective retrieval
  - optimization of conflicting objectives
- Fusion Coder proposed for practical design
- Fusion Coder provides large gains over joint compression (VQ) schemes